NetPoirot: Taking The Blame Game Out of Data Center Operations

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Datacenters can fail.





customer who experienced a nine-day service disruption last month is furious over the software giant's lack of clear and transparent communication about the continuously for two issue.

Netflix

Reddit

problems with two of his foll VMs running in imp' Microsoft's Iowa data center starting April 7. The, that, One of the VMs was down relan the straight days, and was then beh available only ofa The customer, who uses the intermittently for seven exp Azure Virtual Machines days. The other VM was in l Infrastructure-as-a-Service available intermittently for its

Failures are disruptive

- They can cause significant user downtime
 - Loss of revenue for network providers
 - Lower QoE (Quality of Experience) for the users
- This introduces the need for debugging tools

Why is debugging hard?



In the case of a failure...

Someone accepts responsibility

Each blames the other



A real example... Event X

- Azure hypervisors connect to a remote service
- If these connections fail, the VM has uncertain state
 - VM has to reboot
- Did the service fail, or was it the network?

Current tools are insufficient



Can we do better? (Overview)



The monitoring agent

- Runs on all *clients* in our data center
- Captures and digests Windows TCP events
- Reports digests every 30 seconds
- Examples of metrics captured
 - Number of duplicate Acks
 - Number of timeouts
 - Time spent in zero window probing

What is the TCP event digest?

- We aggregate the captured TCP data into epoch digests
- Keep the min, max, 10th,50th,95th percentile, as well as mean and standard deviation across all connections in an epoch for each metric
 - Helps compare performance across the various connections

Why do we think this can work?

- TCP observes the entire communication path
 - It goes through the client, the network, and the server
- It "sees" the failure no matter where it happened
- We know how network failures impact TCP
- How does it react to end point failures?
 - Hard to predict based on protocol design

To distinguish failures...

- We use variants of decision trees
- Other algorithms combine/manipulate features
 - Makes it hard to reason about why they arrive at a decision



Decision trees...

• Greedily pick features that maximize information gain

Will it rain today? His uncertainty is X

Decision trees...

• Greedily pick features that maximize information gain

• Pick the most "informative" features in each step

Will it rain today given that its cloudy outside?



His uncertainty 🖹 X-







Feature 1





What we do to deal with this





Upper portion of an example tree...



What we do to deal with this







Upper portion of an example tree...



95th percentile of the number of timeouts











Other details

- We had to use random forest
 - More stable
- Per application training
- Normalize the data

If thirthrohydrat < x: Send faper data on the same connections

What did we learn from all this?

- TCP sees everything, even at a single end point
 - Allows us to find who was responsible for a failure
- Failures in a group (Client/Server/Network) are similar
 - Makes individual failure classification more challenging
 - Helps NetPoirot be resilient to failures we haven't seen in the past
- The relationship between failures and TCP metrics is non-linear
 - Pearson correlation is low
- Two features suffice to describe each failure as observed by TCP
 - Two largest eigenvalues of the data matrix capture 95% of its variance

Evaluation

- What is the worst case performance?
 - Applications react to failures
 - Their reactions provide useful information
 - But what if this information is not available?
- What if we did not anticipate a failure type?
 - Dormant failures
 - Unknown failures

How did we get labeled data?

- We inject faults into the communication
- Over 6 months of data
- Examples:
 - High CPU load on the client
 - High I/O load on the server
 - Bandwidth throttling in the network
 - Packet reordering in the network

Worse case application

• Only TCP statistics are used from the client side machine



What if we haven't seen the failure before?

	Dormant	Normal		Unkr	iown	
		Client				
High Latency		Server	High Latency			
Reordering		Network	Reordering			
Packet Drop			Packet Drop			
Throttling			Throttling			
High Mem. Client		Higl	n Mem. Client			
High I/O Client		н	igh I/O Client			
High CPU Client	р 	Hi	gh CPU Client			
High Mem. Server	n 	High	Mem. Server			
High I/O Server	р. 	Hi	gh I/O Server			
Slow Reading Server		Slow R	eading Server			
High CPU Server	ř	Hig	h CPU Server			
0) 50% 10	0%	C) 50	%	100%

Performance on real applications

General label	Normal	Client	Networ k			
Precisio n	97.78%	99.7%	100%			
Recall	99.68%	98.25%	99.37			
YouTube						



Event X

Things we did not talk about

- Identifying the actual type of failure
- Sensitivity to machine location
- Aggregation vs per connection classification
- Sensitivity to failure duration
- Modifications to traditional cross validation required

What's next?

- Can we make this application independent?
 - Transfer learning
- Can the end point identify the device causing the failure?
 - Correlate information across clients

Conclusion

- TCP's reactions to network and endpoint failures are significantly different
- We can utilize these differences to find the entity that caused the failure

